

CS 126 Lecture A3: Boolean Logic

Outline

- **Introduction**
- Logic gates
- Boolean algebra
- Implementing gates with switching devices
- Common combinational devices
- Conclusions

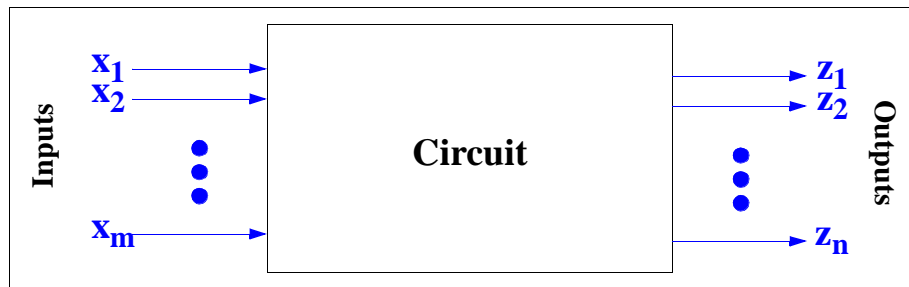
Where We Are At

- We have learned the abstract **interface** presented by a machine: the instruction set architecture
- What we will learn: the **implementation** behind the interface:
 - Start with switching devices (such as transistors)
 - Build logic gates with transistors
 - Build combinational circuit (memory-less) devices using gates
 - Next lecture: build sequential circuit (memory) devices
 - The one after: glue these devices into a computer

Digital Systems

- ... however, the application of digital logic extends way beyond just computers.
- Today, digital systems are replacing all kinds of analog systems in life (data processing, control systems, communications, measurement, ...)
- What is a digital system?
 - Digital: quantities or signals only assume discrete values
 - Analog: quantities or signals can vary continuously
- Why digital systems?
 - Greater accuracy and reliability

Digital Logic Circuits

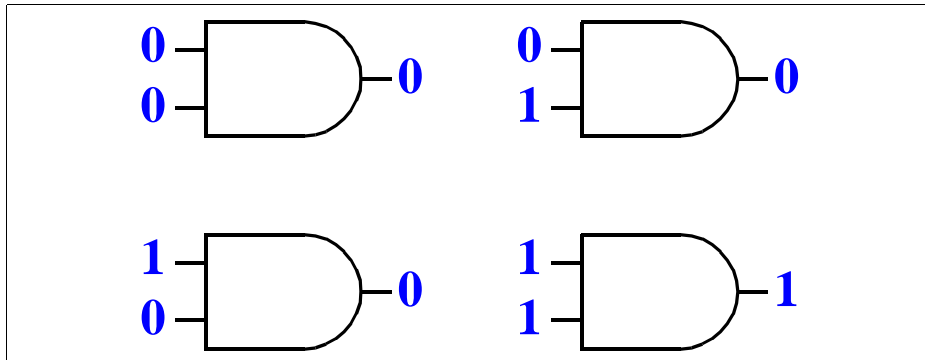


- The heart of a digital system is usually a digital logic circuit

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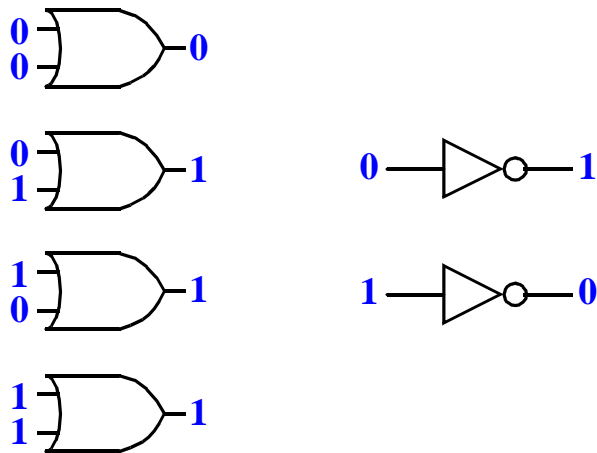
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An AND-Gate

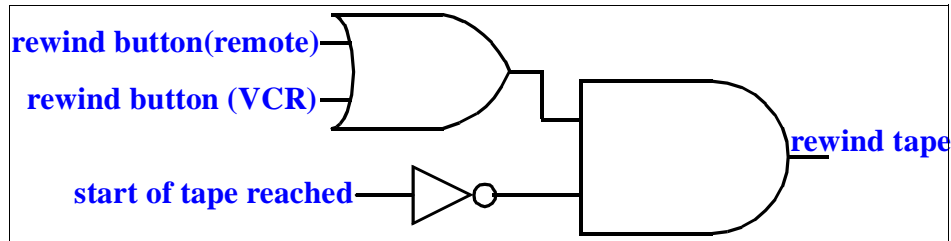


- A smallest useful circuit is a logic gate
- We will connect these small gates into larger circuits

An OR-Gate and a NOT-Gate

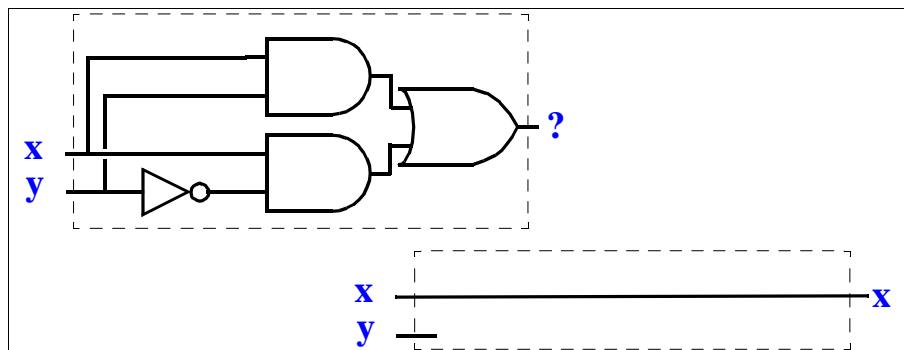


Building Circuits Using Gates



- Can implement **any** circuit using **only** AND, OR, and NOT gates
- But things get complicated when we have lots of inputs and outputs...

Problems



- Many different ways of implementing a circuit (the two above circuits turn out to be the same!)
- How do we find the best implementation? Need better formalism
- Also need more compact representation
- This leads to the study of boolean algebra

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Boolean Algebra

- History
 - Developed in 1847 by Boole to solve mathematic logic problems
 - Shannon first applied it to digital logic circuits in 1939
- Basics
 - **Boolean variables**: variables whose values can be 0 or 1
 - **Boolean functions**: functions whose inputs and outputs are boolean variables
- Relationship with logic circuits
 - Boolean variables correspond to signals
 - Boolean functions correspond to circuits

Defining a Boolean Function with a Truth Table

x	0	0	1	1
y	0	1	0	1
AND(x, y)	0	0	0	1

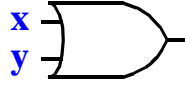
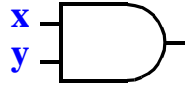
- A systematic way of specifying a function value for **all** possible combination of input values
- A function that takes 2 inputs has 2x2 columns
- A function that takes n inputs has 2ⁿ columns
- This particular example is the AND-function

OR and NOT Truth Tables

x	0	0	1	1
y	0	1	0	1
OR(x, y)	0	1	1	1

x	0	1
NOT(x)	1	0

Defining a General Boolean Function Using Three Basic Boolean Functions



$$\text{AND}(x, y) = xy = x * y$$

$$\text{OR}(x, y) = x + y$$

$$\text{NOT}(x) = x'$$

- The three basic functions have short-hand notations
- Can compose the three basic boolean functions to form arbitrary boolean functions [such as $g(x, y) = xy + z'$]

Two Ways of Defining a Boolean Function

x	0	0	1	1
y	0	1	0	1
XOR(x, y) = x ^ y	0	1	1	0

$$\text{XOR}(x, y) = x \wedge y = x' y + x y'$$

- We have learned that any function can be defined in these two ways: truth table and composition of basic functions
- Why do we need all these different representations?
 - Some are easier than others to begin with to design a circuit
 - Usually start with truth table (or variants of it)
 - Derive a boolean expression from it (perhaps including simplification)
 - Straightforward transformation from boolean expression to circuit

More Examples of Boolean Functions

Sixteen different functions

0 0 1 1	x	
0 1 0 1	y	
0 0 0 0	constant 0	
0 0 0 1	AND (xy)	[decode 11 = 3]
0 0 1 0		[decode 10 = 2]
0 0 1 1	x	
0 1 0 0		[decode 01 = 1]
0 1 0 1	y	
0 1 1 0	XOR (x^y)	
0 1 1 1	OR (x+y)	
1 0 0 0	NOR ("not or")	[decode 00 = 0]
1 0 0 1	== ("not xor")	
1 0 1 0	NOT y (y')	
1 0 1 1		
1 1 0 0	NOT x (x')	
1 1 0 1		
1 1 1 0	NAND ("not and")	
1 1 1 1	constant 1	

Gluing the truth tables of all functions of two variables into one table

For n variables, there are a total of 2^{2^n} functions!

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So How to Translate a Truth Table to a Boolean Expression (Sum-of-Products)?

- form AND terms for each 1 in the function
use v if it corresponds to v = 1
use v' (NOT v) if it corresponds to v = 0
- OR the terms together

Ex: majority function

x:	0	0	0	0	1	1	1	1
y:	0	0	1	1	0	0	1	1
z:	0	1	0	1	0	1	0	1
m:	0	0	0	1	0	1	1	1

$m = x'y'z + xy'z + xyz' + xyz$

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Another Example

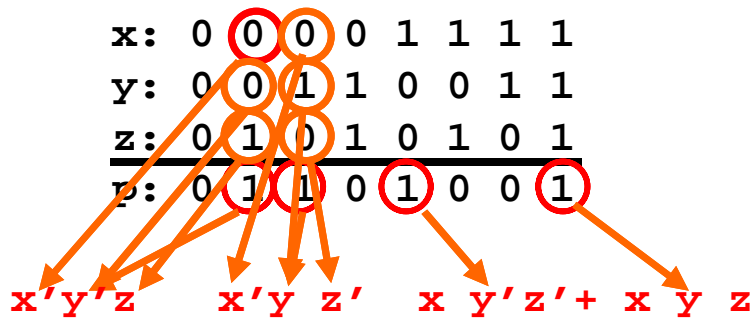
Example: odd parity function

x:	0	0	0	0	1	1	1	1
y:	0	0	1	1	0	0	1	1
z:	0	1	0	1	0	1	0	1
p:	0	1	1	0	1	0	0	1

$$p = x'y'z + x'yz' + xy'z' + xyz$$

001
010
100
111

Parity Function Construction Demo

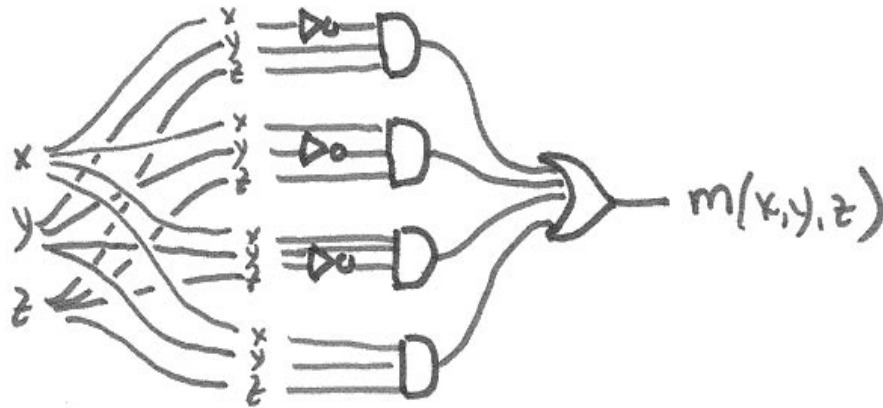


Transform a Boolean Expression into a Boolean Circuit

Use sum-of-products form of function

Example: majority

$$m = x'yz + xy'z + xyz' + xyz$$

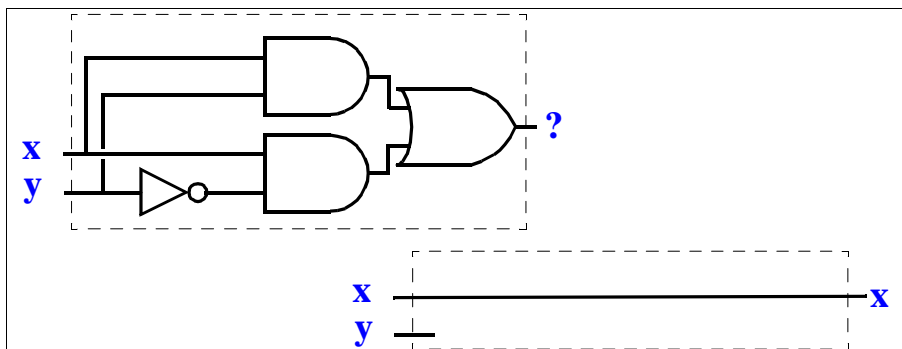


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Simplification Using Boolean Algebra



- Large body of boolean algebra laws can be employed to simplify circuits
- The previous example:

$$xy + xy' = x(y+y') = x*1 = x$$
- Much more, but you don't have to know any of this...

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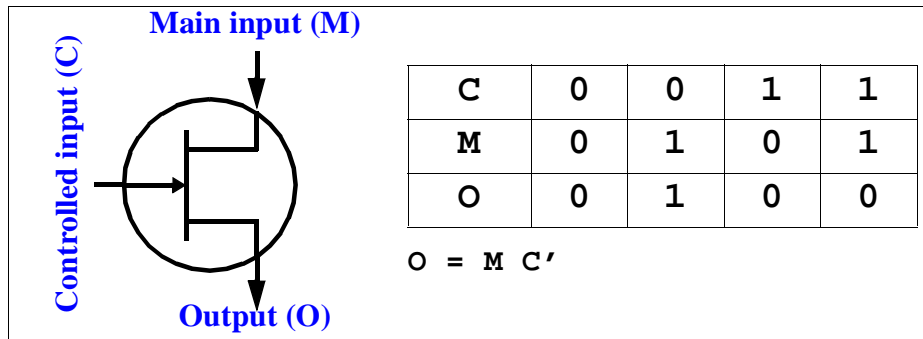
Mini-Summary: How Do We Make a Combinational Circuit

- Represent input signals with input boolean variables, represent output signals with output boolean variables
- Construct truth table based on what we want the circuit to do
- Derive (simplified) boolean expression from the truth table
- Transform boolean expression into a circuit by replacing basic boolean functions with primitive gates

Outline

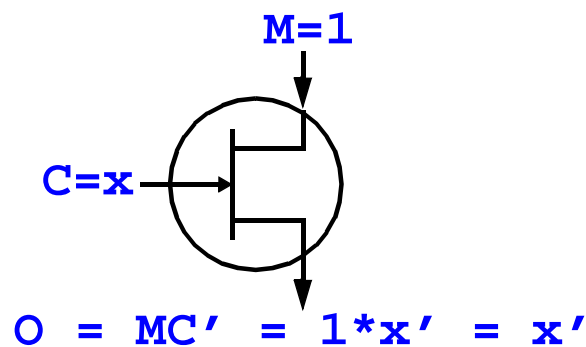
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Switching Devices

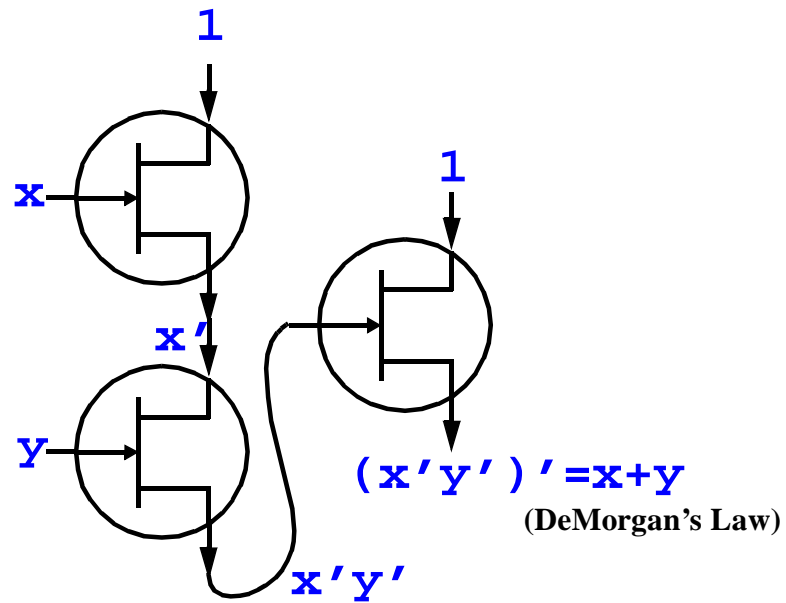


- Any two-state device can be a switching device, examples are relays, diodes, transistors, and magnetic cores
- A transistor example
- Any boolean function can be implemented by wiring together transistors

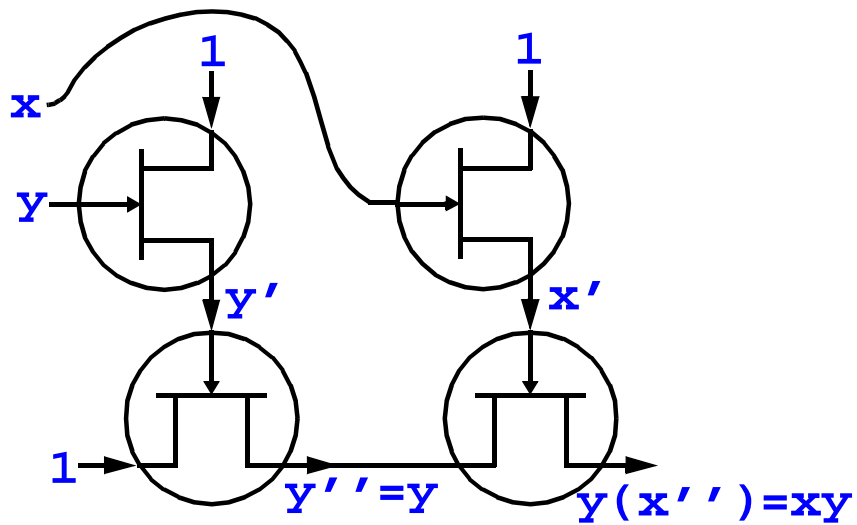
Make a NOT-gate Using a Transistor



Make an OR-gate Using Transistors



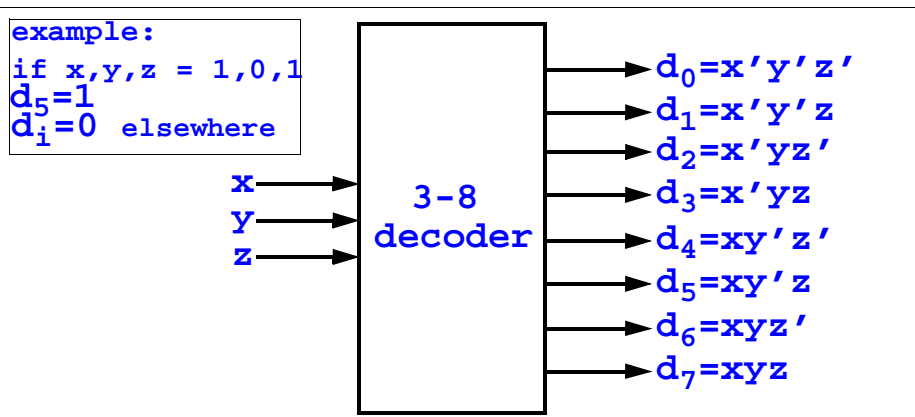
Make an AND-gate Using Transistors



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Decoder Interface



DECODER

N "inputs"

2^N "outputs"

- Turns on precisely one "output"

address is encoded in "inputs"

2^N boolean functions

Deriving Decoder Boolean Expressions

x	0	0	0	0	1	1	1	1
y	0	0	1	1	0	0	1	1
z	0	1	0	1	0	1	0	1
d₀	1	0	0	0	0	0	0	0

$$d_0 = x' y' z'$$

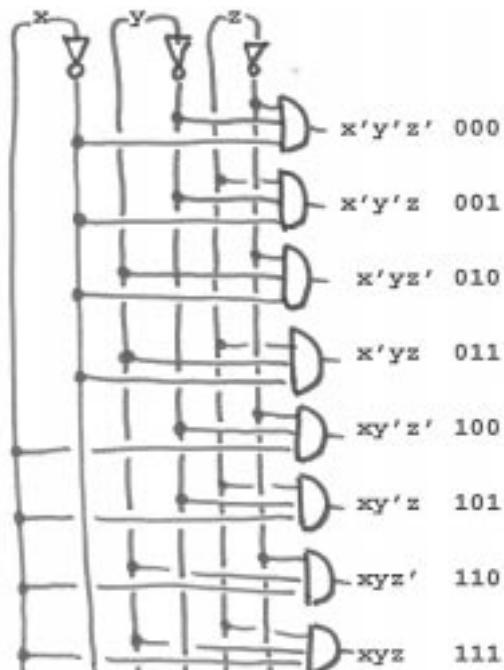
x	0	0	0	0	1	1	1	1
y	0	0	1	1	0	0	1	1
z	0	1	0	1	0	1	0	1
d₁	0	1	0	0	0	0	0	0

$$d_1 = x' y' z$$

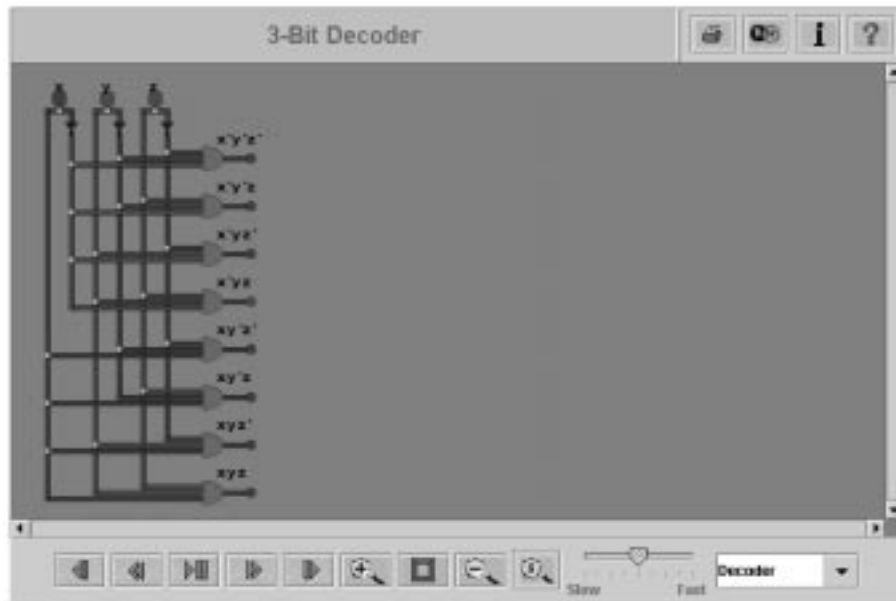
.....

- Can bypass truth table when you're comfortable with this

Decoder Implementation



Decoder Demo

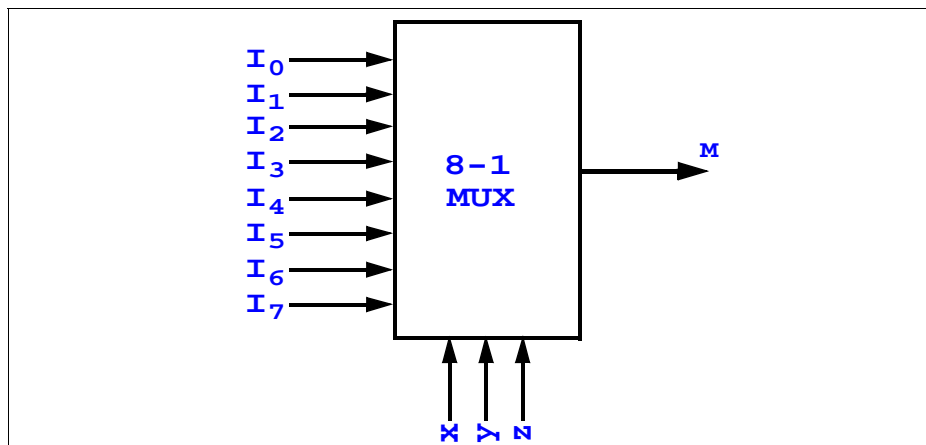


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Multiplexer Interface



- I_0 - I_7 are the “data inputs”, x,y,z form the “control” inputs and are interpreted together as one binary number
- One data input is selected by the control and becomes output
- For example, if x,y,z are 1,0,1, then $M=I_5$

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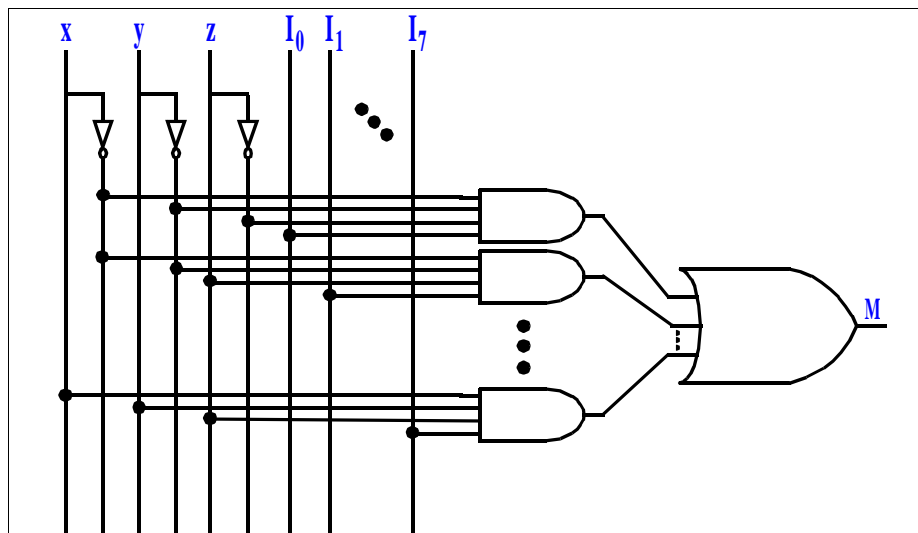
Multiplexer Boolean Expression

x	0	0	0	0	...	1	1
y	0	0	0	0	...	1	1
z	0	0	1	1	...	1	1
I ₇	0	0	0	0	...	0	1
...
I ₁	0	0	0	1	...	0	0
I ₀	0	1	0	0	...	0	0
M	0	1	0	1	...	0	1

$$M = x'y'z'I_0 + x'y'zI_1 + \dots + xyzI_7$$

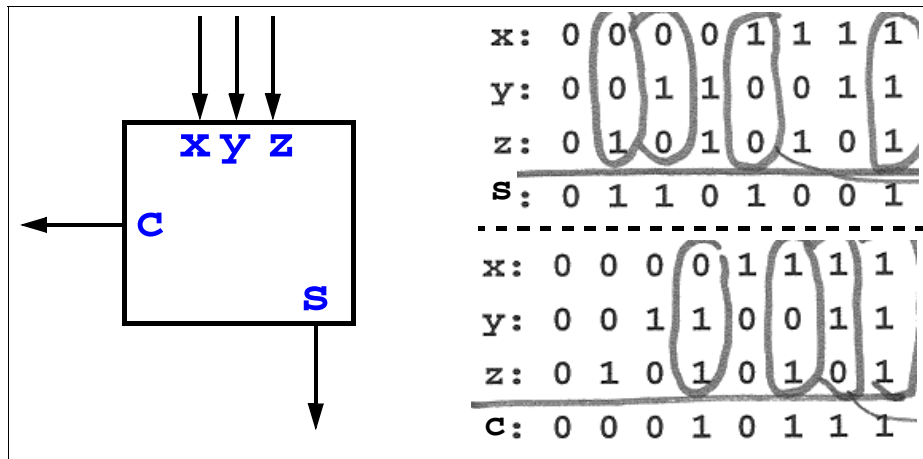
- A lot easier in this case to directly derive the boolean expression instead of starting with a truth table

Multiplexer Implementation



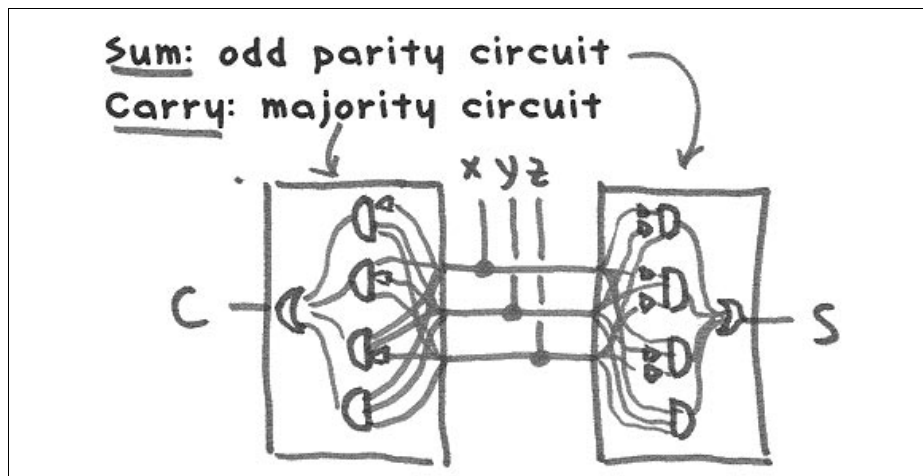
- $M = x'y'z'I_0 + x'y'zI_1 + x'yz'I_2 + x'yzI_3 + xy'z'I_4 + xy'zI_5 + xyz'I_6 + xyzI_7$

An Adder Bit-Slice Interface



- Add three 1-bit numbers x , y , z
- s is the 1-bit sum
- c is the 1-bit carry

An Adder Bit-Slice Implementation



- See slides 11-16, 11-17, and 11-18 for details of the odd parity circuit and majority circuit

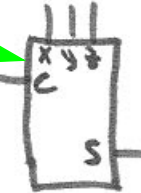
An N-bit Adder Made with Bit-Slices

Goal: add two N-bit numbers x and y

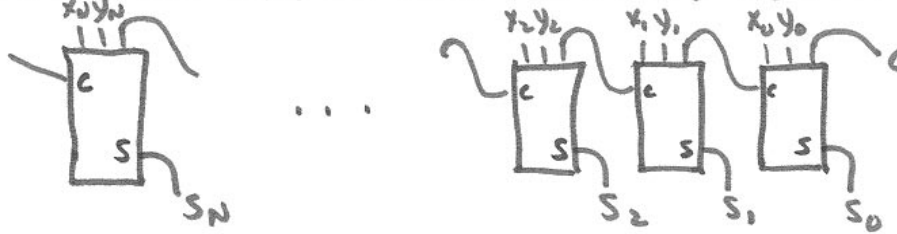
Not a bit-by-bit function of the inputs

because of "carry"

```
  0 1 0 0 1 1 1 1 1
+ 0 0 0 1 0 0 0 0 1
-----
  0 1 1 0 0 0 0 0 0
```



ADDER: one copy for each bit, strung together

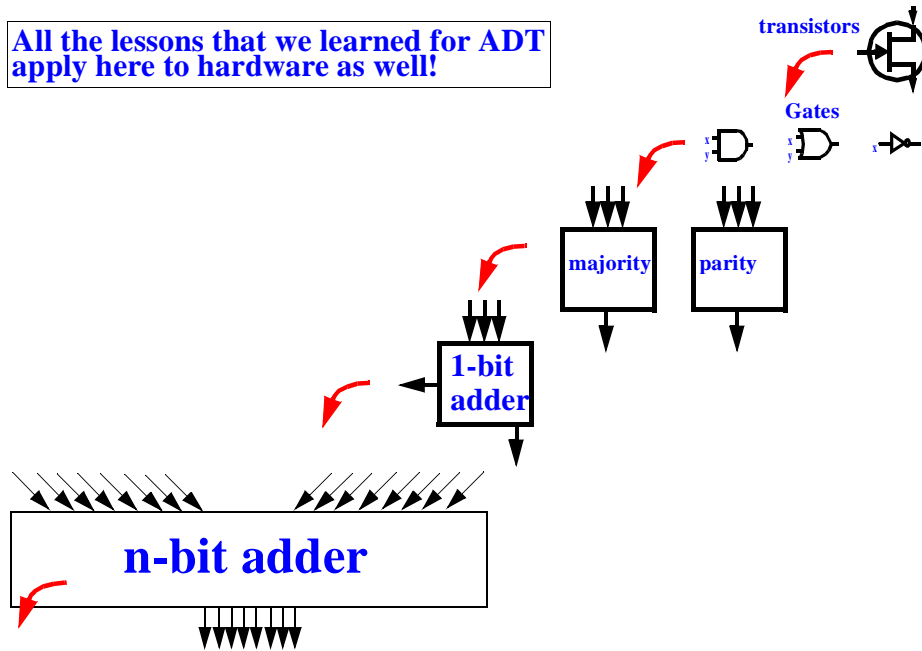


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Abstractions and Encapsulation

All the lessons that we learned for ADT apply here to hardware as well!



Building a Computer Bottom Up

- **Circuit design**: specifying the interconnection of components such as resistors, diodes, and transistors to form logic building blocks
- **Logic design**: determining how to interconnect logic building blocks such as logic gates and flip-flops to form subsystems
- **System design** (or computer architecture): specifying the number, type, and interconnection of subsystems such as memory units, ALUs, and I/O devices

What We Have Learned

- How to build basic gates using transistors
- How to build a combinational circuit
 - Truth table
 - Sum-of-product boolean expression
 - Transform a boolean expression into a circuit of basic gates
- The functionality of some common devices and how they are made
 - Decoder
 - Multiplexer
 - Bit-slice adder
- You're **not** responsible for
 - Boolean algebra laws, or circuit simplification